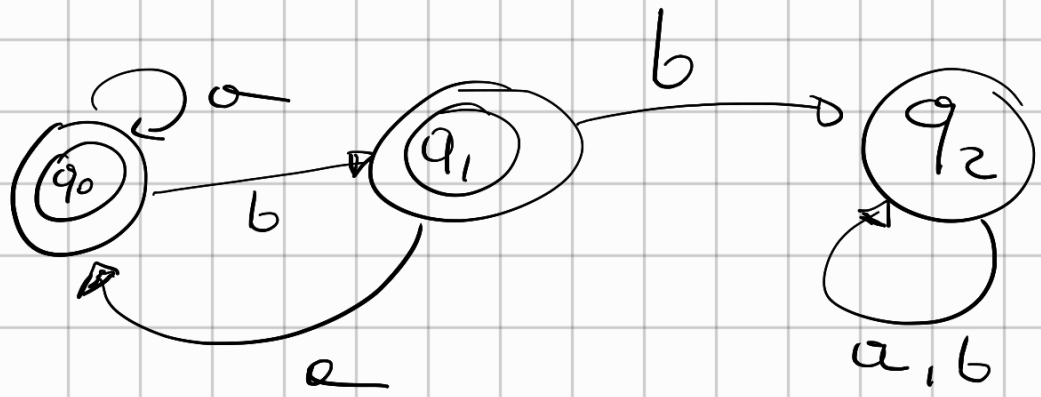


①

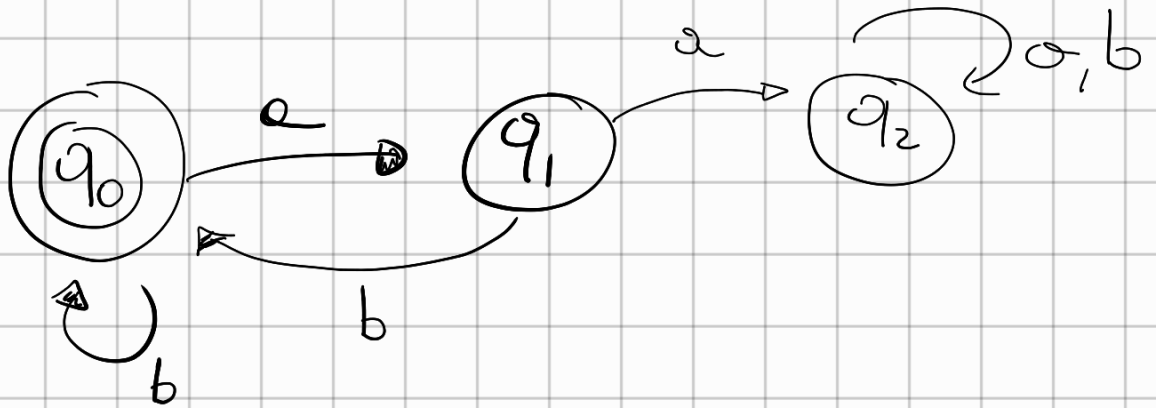


$$(a|ba)^* | (a|ba)^* b \equiv (a|ba)^* (\lambda | b)$$

$$G: NT = \{A, S\} \quad S \rightarrow aS | \lambda | bA$$

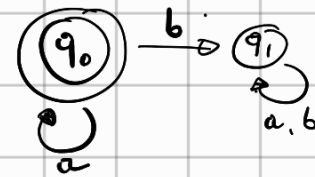
$$T = \{a, b\} \quad A \rightarrow aS | \lambda$$

②



$$\text{Expr: } b^* (a(b)^+)^*$$

$$\begin{cases} S \rightarrow bS | \lambda | aA \\ A \rightarrow bS \end{cases}$$



$$a^+ L = \{a^m \mid m \geq 0\} = \{\lambda, a, a^2, a^3, \dots\}$$



$$L = \{a^m \mid m \geq 0\}$$

$$(a^+)^* \equiv a^*$$

Esempio. abbaab

$$\lambda \quad S \rightarrow aA$$

$aA \quad A \rightarrow bS$

$obS \quad S \rightarrow bS$

$obbS \quad S \rightarrow aA$

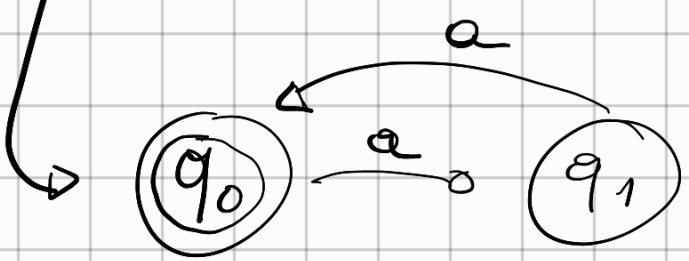
$obbaA \quad A \rightarrow bS$

$obbabS \quad S \rightarrow \lambda$

$obbab$

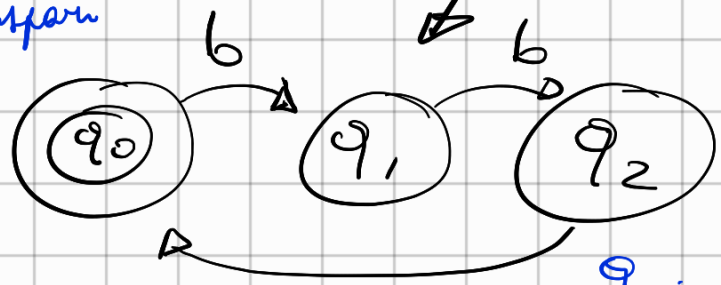
③ → struggle in cui #a pari

$e \quad \underline{\#b \text{ mod } 3 = 0} \quad \leftarrow$



$q_0: a \text{ pari}$

$q_1: a \text{ dispari}$

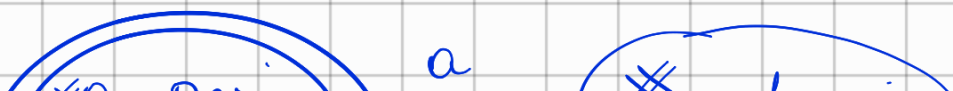


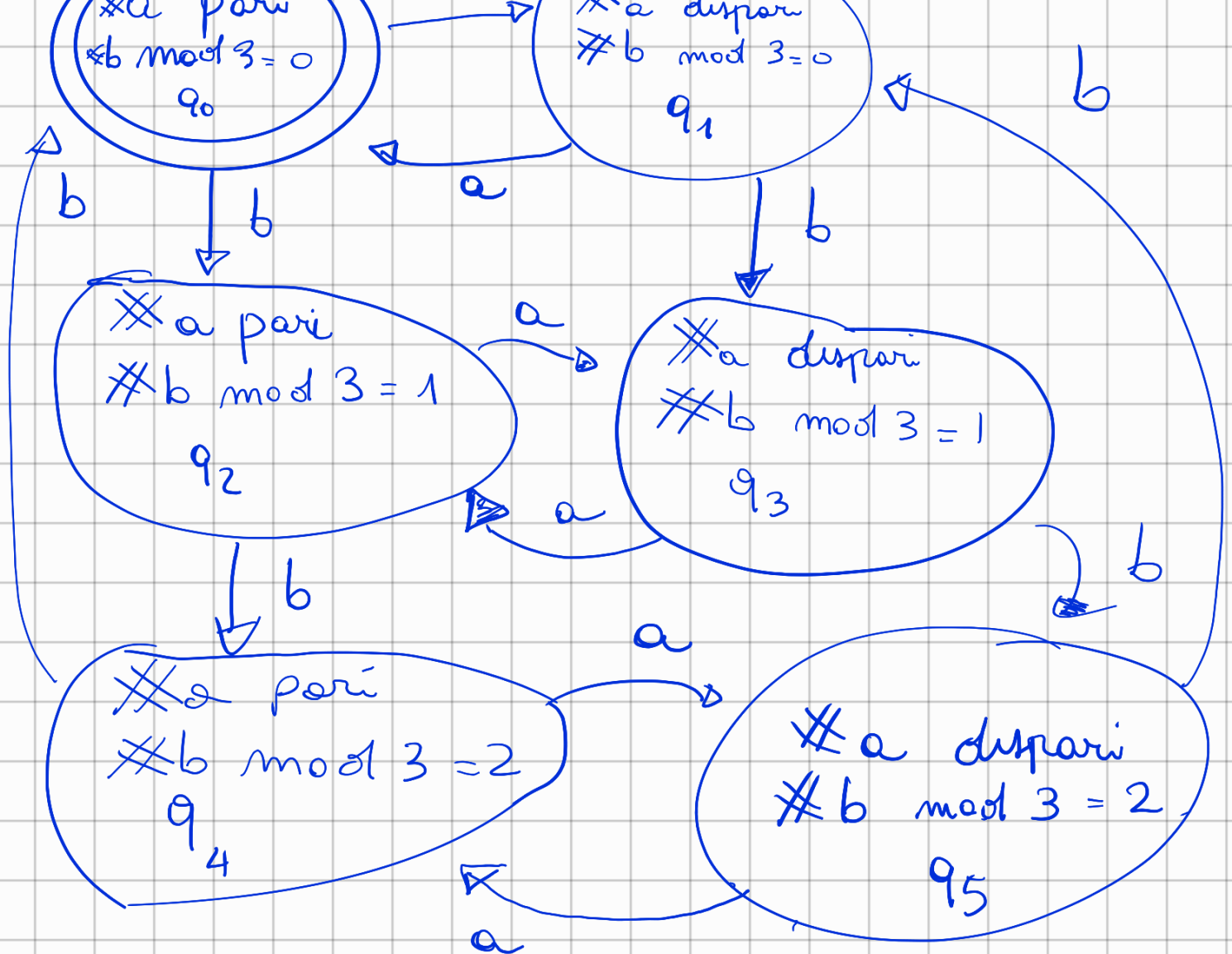
$q_0: \text{mod } 3 = 0$

$q_1: \text{mod } 3 = 1$

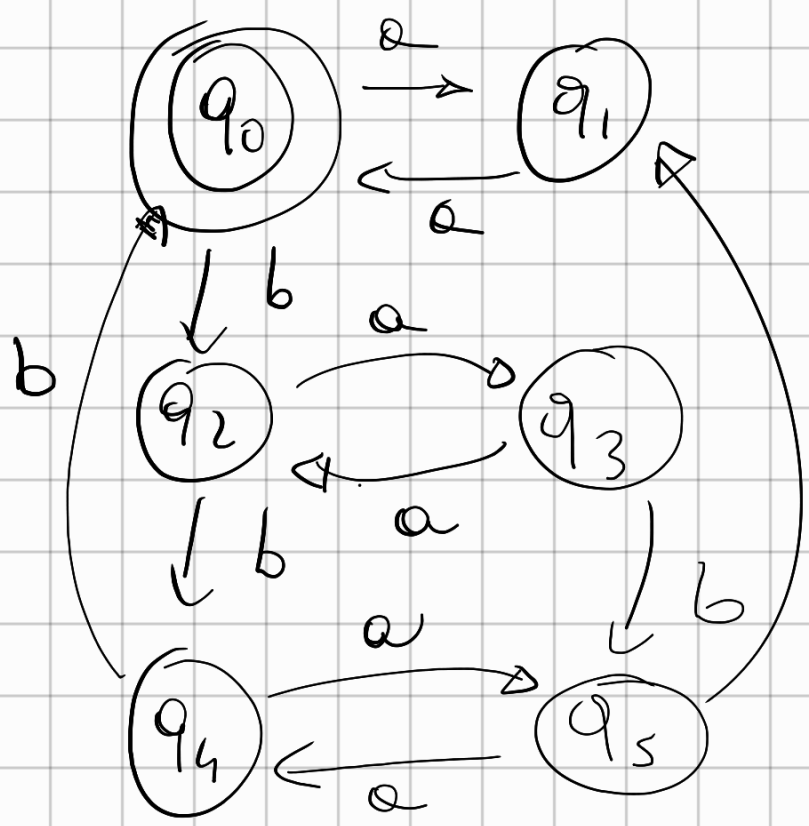
$q_2: \text{mod } 3 = 2$

$a \text{ pari}, \text{ mod } 3 = 0$





Expr: $(aa)^*(bbb)^*$



$\#a$ pari, $\#b \bmod 3 = 0$ $\#a$ dispari, $\#b \bmod 3 = 0$

semantica degli stati

- 1) ~~#a mod 3 = 0, #b dispari~~
 - 2) ~~#a mod 3 = 1, #b dispari~~
 - 3) ~~#a mod 3 = 2, #b dispari~~
 - 4) ~~#a mod 3 = 3, #b dispari~~
 - 4) ~~#a mod 3 = 0, #b pari~~
 - 5) ~~#a mod 3 = 1, #b pari~~
 - 6) ~~#a mod 3 = 2, #b pari~~
 - 8) ~~#a mod 3 = 3, #b pari~~
-

modulo

$$n \text{ mod } r = p$$

$$n \% r = h \text{ resto : } p$$

Esercizi

① costruire un DFA sull'alfabeto $\{a, b\}$ che riconosca le stringhe che contengono aa oppure bb

② Costruire un DFA sull'alfabeto $\{a, b\}$
che riconosca le stringhe che
contengono aa e bb

③ Costruire un DFA sull'alfabeto $\{a, b\}$
che riconosca le stringhe che
iniziano con ab e finiscono con ba